

Preconditioning for Stochastic Automata Networks

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Outline

- SANs
- Literature Review
- NKP Preconditioner: applied to general matrix, multilinear algebra
- NKP Preconditioner for SAN's Q matrix in descriptor form
- Implementation issues
- Suggestions

Markov chains

- a stochastic process which follows the Markov property. The state of the system at time t only depends on the most recent past.

- DTMC

$$P = \begin{pmatrix} p_{1,1} & p_{1,2} & \cdots & p_{1,n} \\ p_{2,1} & p_{2,2} & \cdots & p_{2,n} \\ \vdots & \vdots & \ddots & \vdots \\ p_{n,1} & p_{n,2} & \cdots & p_{n,n} \end{pmatrix},$$

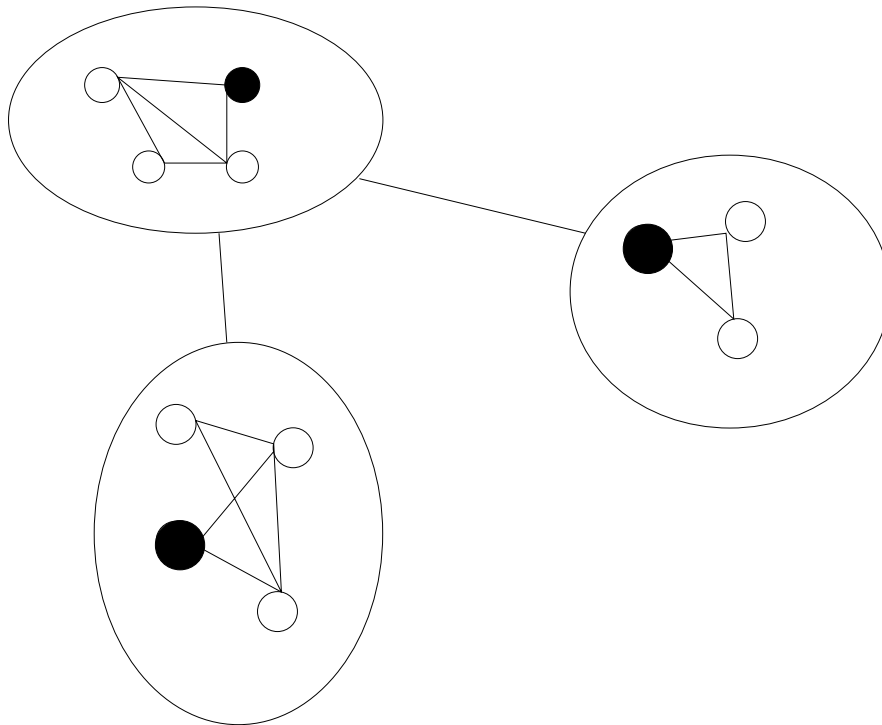
where $p_{i,j}$ holds the conditional probability of moving from state i to state j in one time step.

- CTMC

$$Q = \begin{pmatrix} q_{1,1} & q_{1,2} & \cdots & q_{1,n} \\ q_{2,1} & q_{2,2} & \cdots & q_{2,n} \\ \vdots & \vdots & \ddots & \vdots \\ q_{n,1} & q_{n,2} & \cdots & q_{n,n} \end{pmatrix},$$

where $q_{i,j}$ holds the transition rate of moving from state i to state j

and $q_{i,i} = -\sum_{j \neq i} q_{i,j}$.



Stochastic Automata Networks (SANs)

- Problem with MC analysis = size of P or Q .
- One solution = SANs (represent Q in compact form)
- A SAN is a collection of stochastic automata which act more or less independently requiring only infrequent interaction.

Tensor Algebra: \otimes, \oplus are key operations in SAN formalism.

- $A \in \mathfrak{R}^{m_1 \times n_1}, B \in \mathfrak{R}^{m_2 \times n_2}$, then $A \otimes B \in \mathfrak{R}^{m_1 m_2 \times n_1 n_2}$ and

is defined by

$$A \otimes B = \begin{pmatrix} a_{1,1}B & \cdots & a_{1,n_1}B \\ \vdots & \ddots & \vdots \\ a_{m_1,1}B & \cdots & a_{m_1,n_1}B \end{pmatrix}.$$

- $A \oplus B = A \otimes I_{m_2} + I_{m_1} \otimes B$. (\oplus is defined in terms of \otimes .)

(defined only for square matrices)

- Some Tensor Properties:

$$- (A \otimes B)^{-1} = A^{-1} \otimes B^{-1}.$$

$$- \bigoplus_{i=1}^N A^{(i)} = \sum_{i=1}^N I_{n_1} \otimes \cdots \otimes I_{n_{i-1}} \otimes A^{(i)} \otimes I_{n_{i+1}} \otimes \cdots \otimes I_{n_N}.$$

SAN descriptor

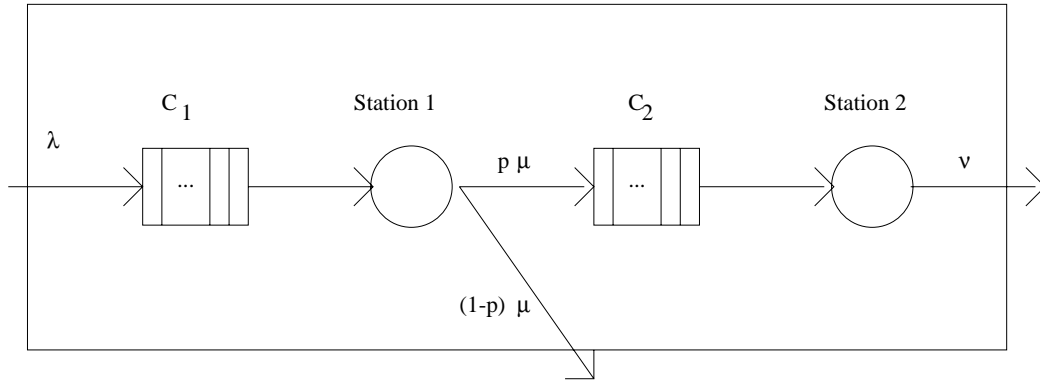
- MC has global Q which can be stored in much more compact form as the ordinary sum of Kronecker products of much smaller matrices.

$$Q = \sum_{j=1}^{2E+N} \otimes_{i=1}^N Q_j^{(i)},$$

where

N = number of automata,

E = number of synchronizing events.



SAN Example: A Simple Queueing Network

Consider a queueing network with two exponential, finite-capacity, single-server stations in tandem. When a station is full, customers are lost. For example, for $C_1 = 1$ and $C_2 = 2$, the transition rate matrix is given by

$$Q = \begin{pmatrix} -\lambda & 0 & 0 & \lambda & 0 & 0 \\ \nu & -(\nu + \lambda) & 0 & 0 & \lambda & 0 \\ 0 & \nu & -(\nu + \lambda) & 0 & 0 & \lambda \\ \mu(1-p) & \mu p & 0 & -\mu & 0 & 0 \\ 0 & \mu(1-p) & \mu p & \nu & -(\mu + \nu) & 0 \\ 0 & 0 & \mu(1-p) & 0 & \nu & -(\nu + \mu(1-p)) \end{pmatrix}.$$

Writing Q as a SAN

- In general, a SAN represents the global generator matrix as

$$Q = \bigoplus_{i=1}^N Q_l^{(i)} + \sum_{j=1}^E Q_{e_j},$$

where

N = number of automata,

E = number of synchronizing events,

$Q_l^{(i)}$ = automata i 's local transition matrix,

Q_{e_j} = synchronizing event j 's contribution to the global Q .

SAN Example: Writing Q as a SAN

- Our example has only 1 synchronizing event; a customer leaves station 1 and enters station 2. $A^{(1)}$ = station 1's automata and $A^{(2)}$ = station 2's automata. Thus,

$$\begin{aligned}
 Q &= Q_l^{(1)} \oplus Q_l^{(2)} + Q_{e_1} \\
 &= \begin{pmatrix} \lambda & & \lambda \\ \mu(1-p) & & -\mu(1-p) \end{pmatrix} \oplus \begin{pmatrix} 0 & 0 & 0 \\ \nu & -\nu & 0 \\ 0 & \nu & -\nu \end{pmatrix} + \\
 &\begin{pmatrix} 0 & 0 \\ 1 & 0 \end{pmatrix} \otimes \begin{pmatrix} 0 & \mu p & 0 \\ 0 & 0 & \mu p \\ 0 & 0 & 0 \end{pmatrix} + \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix} \otimes \begin{pmatrix} -\mu p & 0 & 0 \\ 0 & -\mu p & 0 \\ 0 & 0 & 0 \end{pmatrix}.
 \end{aligned}$$

- We only need to store $Q_l^{(1)}$, $Q_l^{(2)}$, and the 4 matrices making up Q_{e_1} .

For very large MCs, this is a huge savings in storage. The Global Q never needs to be formed or stored.

Stationary analysis of a SAN

GOAL: Solve the following system for π

$$\pi \left(\sum_{j=1}^{2E+N} \otimes_{i=1}^N Q_j^{(i)} \right) = 0,$$
$$\pi e = 1.$$

- Review of methods employed
 - Direct: L, U factors hard to get
 - Iterative: Power–Stewart. Splittings–Dayar.
 - Projection: Arnoldi, GMRES–Stewart. BiCGSTAB, QMR–Buchholz.
 - Decompositional: hierarchical SANS–Buchholz, NCD SANS–Dayar.
- Problem: SANS \downarrow storage, \uparrow complexity \Rightarrow time til convergence=high.
- Solutions:
 - RSS - Ciardo, Miner, Kemper
 - Equivalent SANS - Buchholz
 - Preconditioners - Stewart, Buchholz, Chan

Existing SAN Preconditioners

- *ILU* Preconditioners are hard to employ.
- Stewart's Neumann series inverse: M^{-1} approximates $Q^\#$.

$$M^{-1} = \sum_{m=0}^M P^m.$$

- Buchholz's preconditioner, similar to Stewart's.
- Additive/Multiplicative Schwartz
- Diagonal: $M^{-1} = D^{-1}$, where D is the diagonal of Q .

Potential SAN Preconditioner: NKP ???

NKP Preconditioner for general matrix R

- Find A, B so that $R \approx A \otimes B$.

Precondition with $M^{-1} = A^{-1} \otimes B^{-1}$.

- Pitsianis and Van Loan find A, B so that $\|R - A \otimes B\|_F^2$ is minimized by transforming this problem to an equivalent, easier problem.
- They show $\|R - A \otimes B\|_F^2 = \|\tilde{R} - ab^T\|_F^2$, where \tilde{R} is a special rearrangement of the matrix R , $a = \text{vec}(A)$ and $b = \text{vec}(B)$.

Solving NKP Problem: $\min_{a,b} \|\tilde{R} - ab^T\|_F^2$.

- Rank-1 approximation problem has SVD solution.

$$a^* = \sigma_1 u_1, b^* = v_1.$$

- Reversing the vectorizing operation gives the optimal A, B matrices.

Advantages of NKP Preconditioner:

- To find A, B , only need \tilde{R} -vector multiplication algorithm.

\tilde{R} does not need to be formed or stored.

- Only little matrices A^{-1}, B^{-1} need to be formed and stored.

Efficient vector-Kronecker product multiplication algorithm can be used so that M^{-1} is not explicitly needed.

The NKP, $A \otimes B$, when $R = \sum_{i=1}^p G_i \otimes F_i$

- $\tilde{R} = \sum_{i=1}^p g_i f_i^T$ where $g_i = \text{vec}(G_i)$ and $f_i = \text{vec}(F_i)$.
- It can be proven that when R has this special structure, the optimal A, B matrices are linear combinations of the matrices making up R .

$$A = \alpha_1 G_1 + \alpha_2 G_2 + \cdots + \alpha_p G_p,$$

$$B = \beta_1 F_1 + \beta_2 F_2 + \cdots + \beta_p F_p.$$

- Original NKP problem can be written as a nonlinear optimization problem. Choose $\alpha_1, \alpha_2, \dots, \alpha_p, \beta_1, \beta_2, \dots, \beta_p$ so that this nonlinear convex function is minimized

$$\begin{aligned} \|R - A \otimes B\|_F^2 &= \left\| \sum_{i=1}^p G_i \otimes F_i - \left(\sum_{i=1}^p \alpha_i G_i \right) \otimes \left(\sum_{i=1}^p \beta_i F_i \right) \right\|_F^2 \\ &= \sum_{i=1}^{n_G} \sum_{j=1}^{n_F} \left[\left(\sum_{k=1}^p \sigma_{G_k}(i) \cdot \sigma_{F_k}(j) \right) - \left(\sum_{k=1}^p \alpha_k \sigma_{G_k}(i) \right) \left(\sum_{k=1}^p \beta_k \sigma_{F_k}(i) \right) \right]^2. \end{aligned}$$

Back to SANs

- $Q = \sum_{j=1}^{2E+N} \otimes_{i=1}^N Q_j^{(i)}$ has special structure similar to $R = \sum_{i=1}^p G_i \otimes F_i$.

But Q is Kronecker product of N terms, not just 2 terms.

- SAN Problem: find A, B, \dots, N to min $\|Q - A \otimes B \otimes \dots \otimes N\|_F^2$.

Precondition with $M^{-1} = A^{-1} \otimes B^{-1} \otimes \dots \otimes N^{-1}$.

- Questions:

– Finding A, B was based on the matrix SVD.

Extension for finding A, B, \dots, N ?

– When $R = \sum_{i=1}^p G_i \otimes F_i$, then $A = \sum_{i=1}^p \alpha_i G_i$, $B = \sum_{i=1}^p \beta_i F_i$.

Extension to case when $R = Q = \sum_{j=1}^{2E+N} \otimes_{i=1}^N Q_j^{(i)}$?

– To find A, B , we rearranged R to form \tilde{R} . For A, B, \dots, N , what

is the proper rearrangement of Q ?

Multilinear Algebra

- The tensor is fundamental object of multilinear algebra.
 - 1^{st} -order tensor is a vector.
 - 2^{nd} -order tensor is a matrix.
 - 3^{rd} -order tensor is a 3D box.
- Operations on Tensors:
 - inner product
 - outer product
 - scalar product
 - tensor-matrix multiplication
- Outer product of N vectors $(a \circ b \circ \dots \circ n)$ results in an N^{th} -order tensor.

The HOSVD

- Tensors have rank as well as a SVD, called the HOSVD.
- A rank-1 N^{th} -order tensor is written as the outer product of N vectors.
- HOSVD of N^{th} -order tensor R shows that R can be written as the sum of rank-1 tensors.

$$R = \sum_{i_1=1}^{I_1} \sum_{i_2=1}^{I_2} \cdots \sum_{i_N=1}^{I_N} s_{i_1, i_2, \dots, i_N} u_{i_1}^{(1)} \circ u_{i_2}^{(2)} \circ \cdots \circ u_{i_N}^{(N)}.$$

Multilinear Algebra and the NKP

- We want to min $\|R - A \otimes B \otimes \cdots \otimes N\|_F^2$ for a general matrix R .
- Following the 2D case, we need to show $\exists \tilde{R} \ni$

$$\|R - A \otimes B \otimes \cdots \otimes N\|_F^2 = \|\tilde{R} - a \circ b \circ \cdots \circ n\|_F^2.$$

- $a \circ b \circ \cdots \circ n$ is a rank-1 N^{th} -order tensor
 \Rightarrow rearrangement of R (\tilde{R}) must be N^{th} -order tensor.
- Goal: approximate N^{th} -order tensor with rank-1 N^{th} -order tensor
 \Rightarrow use truncated HOSVD.

Question: How do we define this N^{th} -order tensor \tilde{R} ?

Defining Rearrangement Operator for Higher Orders

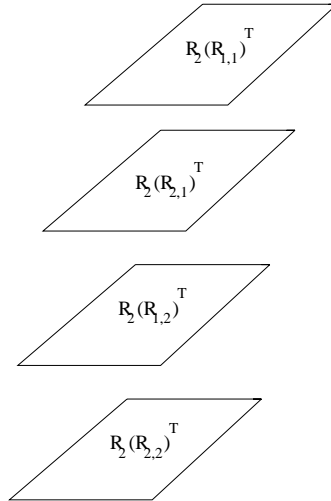
- For 2^{nd} -order tensor R , Pitsianis and Van Loan have already defined $\tilde{R} = R_2(R)$.
- The 3^{rd} -order rearrangement operator, $R_3(R)$, uses the 2^{nd} -order rearrangement operator.
- In general, we define the rearrangement operator recursively.

We can then prove that with this recursive definition, the sums of squares match.

Example: Approximate $R_{12 \times 12}$ by $A_{2 \times 2}, B_{2 \times 2}, C_{3 \times 3}$.

$$R_{12 \times 12} = \begin{pmatrix} R_{1,1} & R_{1,2} \\ R_{2,1} & R_{2,2} \end{pmatrix}.$$

$$\tilde{R}_{4 \times 4 \times 9} = R_3(R) =$$



and then

$$\|R - A \otimes B \otimes C\|_F^2 = \|R_3(R) - a \circ b \circ c\|_F^2.$$

Approximating N^{th} -order tensor by a rank-1

N^{th} -order tensor

We want to min $\|\tilde{R} - a \circ b \circ \dots \circ n\|_F^2$. where \tilde{R} is an N^{th} -order tensor and $a \circ b \circ \dots \circ n$ is a rank-1 N^{th} -order tensor.

- Unfortunately, truncating the HOSVD does not give the optimal rank-1 approximation to the N^{th} -order tensor \tilde{R} .
- de Lathauwer:
 - truncated HOSVD is *good* approximation.
 - HO power algorithm to find *optimal* rank-1 approximation.
- We are only looking for an approximate inverse preconditioner
 \Rightarrow try truncating the HOSVD. Let

$$a = s_{1,1,\dots,1} u_1^{(1)}, b = u_1^{(2)}, \dots, n = u_1^{(N)}.$$

We hope $A \otimes B \otimes \dots \otimes N$ is a good approximation to R .

Practical Issues

- We'd rather not form the N^{th} -order tensor \tilde{R} or find its HOSVD.

That'd be counter to the purpose of using a SAN.

- \Rightarrow try to extend Pitsianis and Van Loan's special structure work.

- Recall: $R = \sum_{i=1}^p G_i \otimes F_i,$

$$\Rightarrow \text{optimal } A = \sum_{i=1}^p \alpha_i G_i, \text{ optimal } B = \sum_{i=1}^p \beta_i F_i.$$

- NKP Problem transformed into a nonlinear convex optimization problem, where \tilde{R} and the HOSVD of \tilde{R} were never needed.

Question: Does this extend to Q 's special structure?

Solving $\| \sum_{j=1}^T \otimes_{i=1}^N Q_j^{(i)} - A \otimes B \otimes \cdots \otimes N \|_F^2$, **where**

$$T = 2E + N.$$

- We can show

$$A \approx \alpha_1 Q_1^{(1)} + \alpha_2 Q_2^{(1)} + \cdots + \alpha_T Q_T^{(1)}$$

$$B \approx \beta_1 Q_1^{(2)} + \beta_2 Q_2^{(2)} + \cdots + \beta_T Q_T^{(2)}$$

⋮

$$N \approx \eta_1 Q_1^{(N)} + \eta_2 Q_2^{(N)} + \cdots + \eta_T Q_T^{(N)}.$$

- Thus, just as in the optimal 2-matrix case, the original problem

$$\min \|Q - A \otimes B \otimes \cdots \otimes N\|_F^2$$

can be transformed into a nonlinear convex optimization problem.

$$\|Q - A \otimes B \otimes \cdots \otimes N\|_F^2 \approx \left\| \sum_{j=1}^T \otimes_{i=1}^N Q_j^{(i)} - \left(\sum_{j=1}^T \alpha_j Q_j^{(1)} \right) \otimes \right. \\ \left. \left(\sum_{j=1}^T \beta_j Q_j^{(2)} \right) \otimes \cdots \otimes \left(\sum_{j=1}^T \eta_j Q_j^{(N)} \right) \right\|_F^2$$

$$= \sum_{a=1}^{n_{Q(1)}} \sum_{b=1}^{n_{Q(2)}} \cdots \sum_{n=1}^{n_{Q(N)}} \left[\left(\sum_{p=1}^T \sigma_{Q_p^{(1)}}(a) \cdot \sigma_{Q_p^{(2)}}(b) \cdots \sigma_{Q_p^{(N)}}(n) \right) - \right. \\ \left. \left(\sum_{p=1}^T \alpha_p \sigma_{Q_p^{(1)}}(a) \right) \left(\sum_{p=1}^T \beta_p \sigma_{Q_p^{(2)}}(b) \right) \cdots \left(\sum_{p=1}^T \eta_p \sigma_{Q_p^{(N)}}(n) \right)^2 \right]$$

With this transformation, \tilde{R} and its HOSVD are never needed.

- This is a nonlinear convex function of NT variables.

As N and $E \uparrow$, practicality of problem transformation \downarrow .

But, we can group automata!

Open Issues

- What values of N, T make transforming the preconditioning problem into a nonlinear optimization problem worthwhile?
- Is there a way to do a vector- \tilde{R} multiplication without forming \tilde{R} or the global Q ? If so, we do not have to resort to the nonlinear optimization problem. We could use de Lathauwer's best rank-1 approximation algorithm to find the *optimal* A, B, \dots, N matrices.
- Does $(A \otimes B \otimes \dots \otimes N \approx Q) \Rightarrow (A^{-1} \otimes B^{-1} \otimes \dots \otimes N^{-1} \approx Q^{-1})$?

The End